Lecture 7 Visualization, Randomness, and Computation

History of Data Science, Spring 2022 @ UC San Diego Suraj Rampure

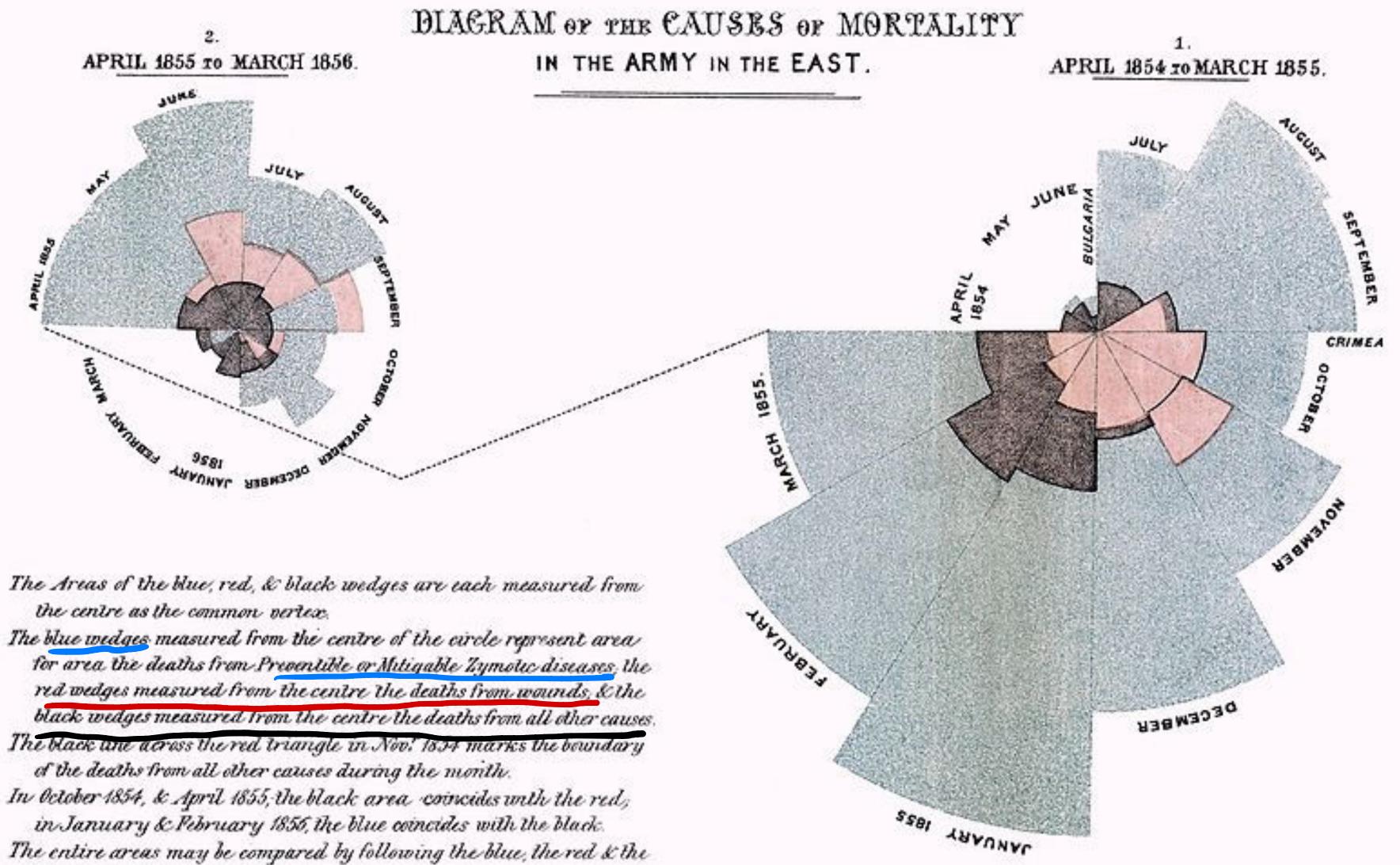
Announcements

- Look at others' Homework 6 websites!
 - Posted in the #homework6 channel on Slack.
 - If your images don't load, please fix them ASAP, and message if you need help.
- Slack • Homework 7 will be released by **noon tomorrow**. Make sure to read Homework 5 solutions (posted on Campuswire)!
- - Many misconceptions about the Law of Large Numbers.



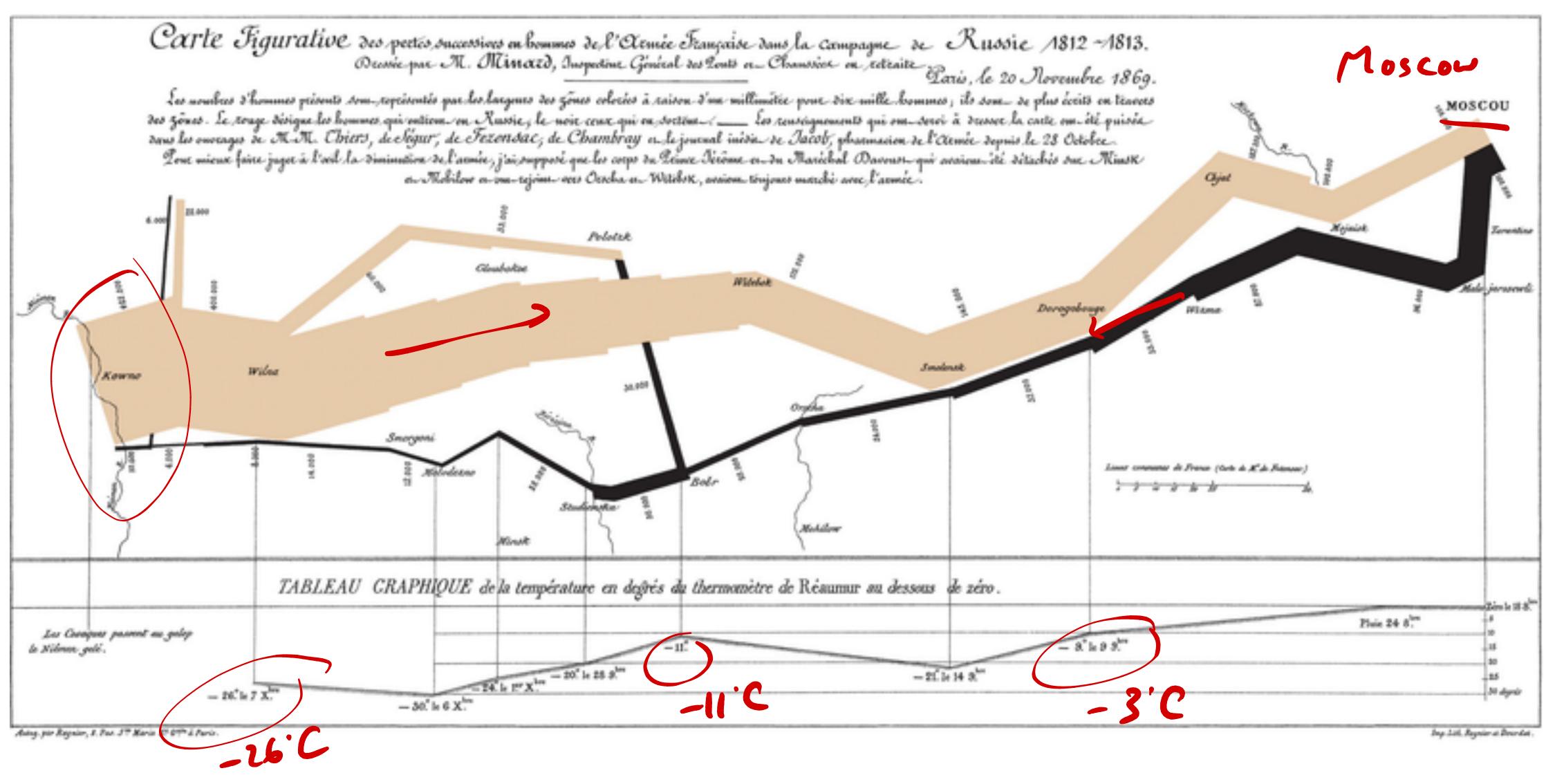
- Finish visualization.
- Revisit the development of the Gaussian distribution.
- Abacus and the binary number system.

Visualization



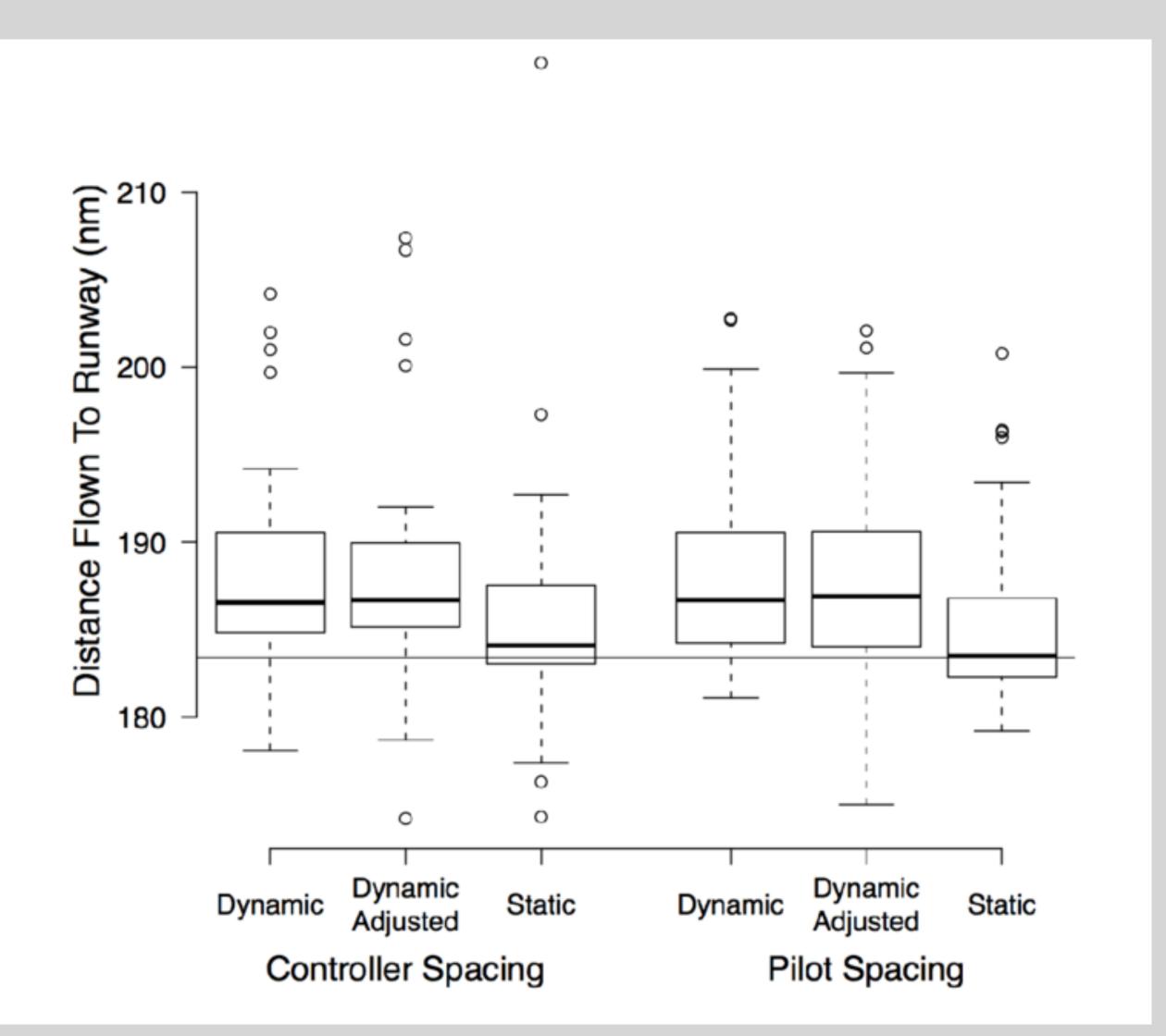
- The entire areas may be compared by following the blue, the red & the black lines enclosing them.

1855: Florence Nightingale's depiction of the deaths of British soldiers in the Crimean war. Florence Nightingale is known as the founder of modern nursing.



1869: Charles Joseph Minard's visualization of the French invasion of Russia (led by Napoleon).

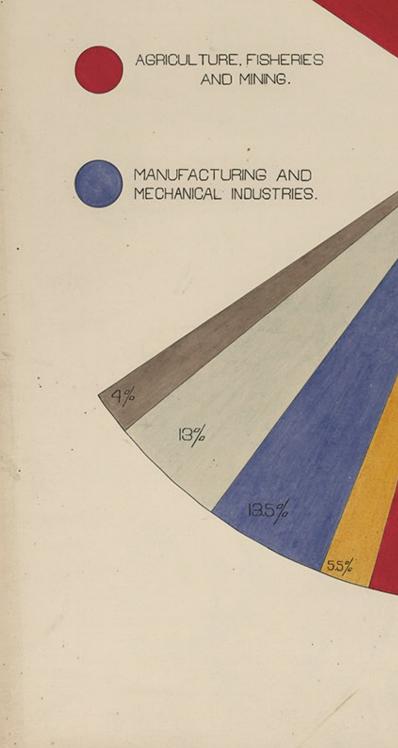


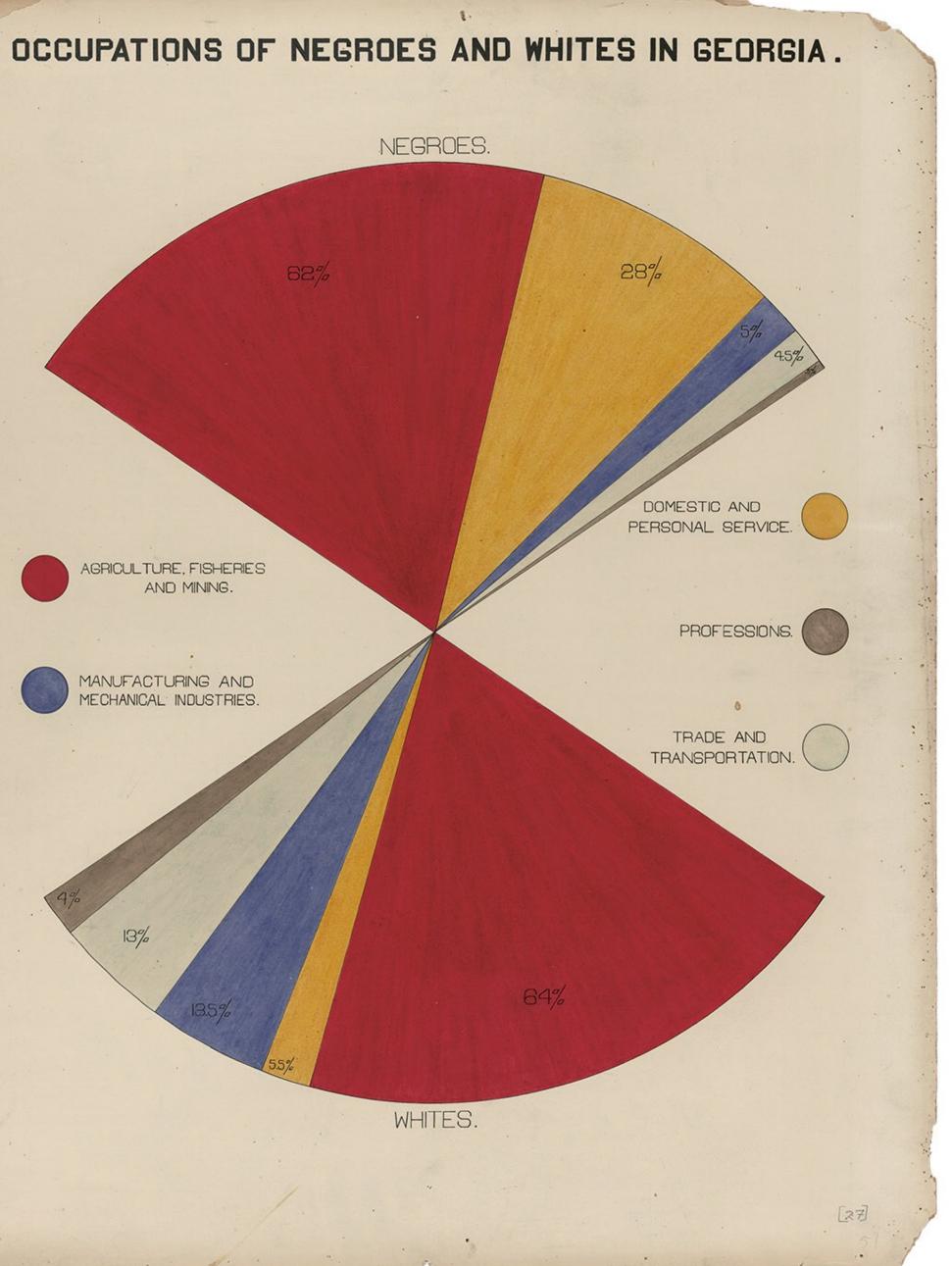


1973: John Tukey, who defined the term "Exploratory Data Analysis", created the box plot, which describes a numerical distribution using a 5 number summary.

W. E. B. Dubois' visualizations of Black America

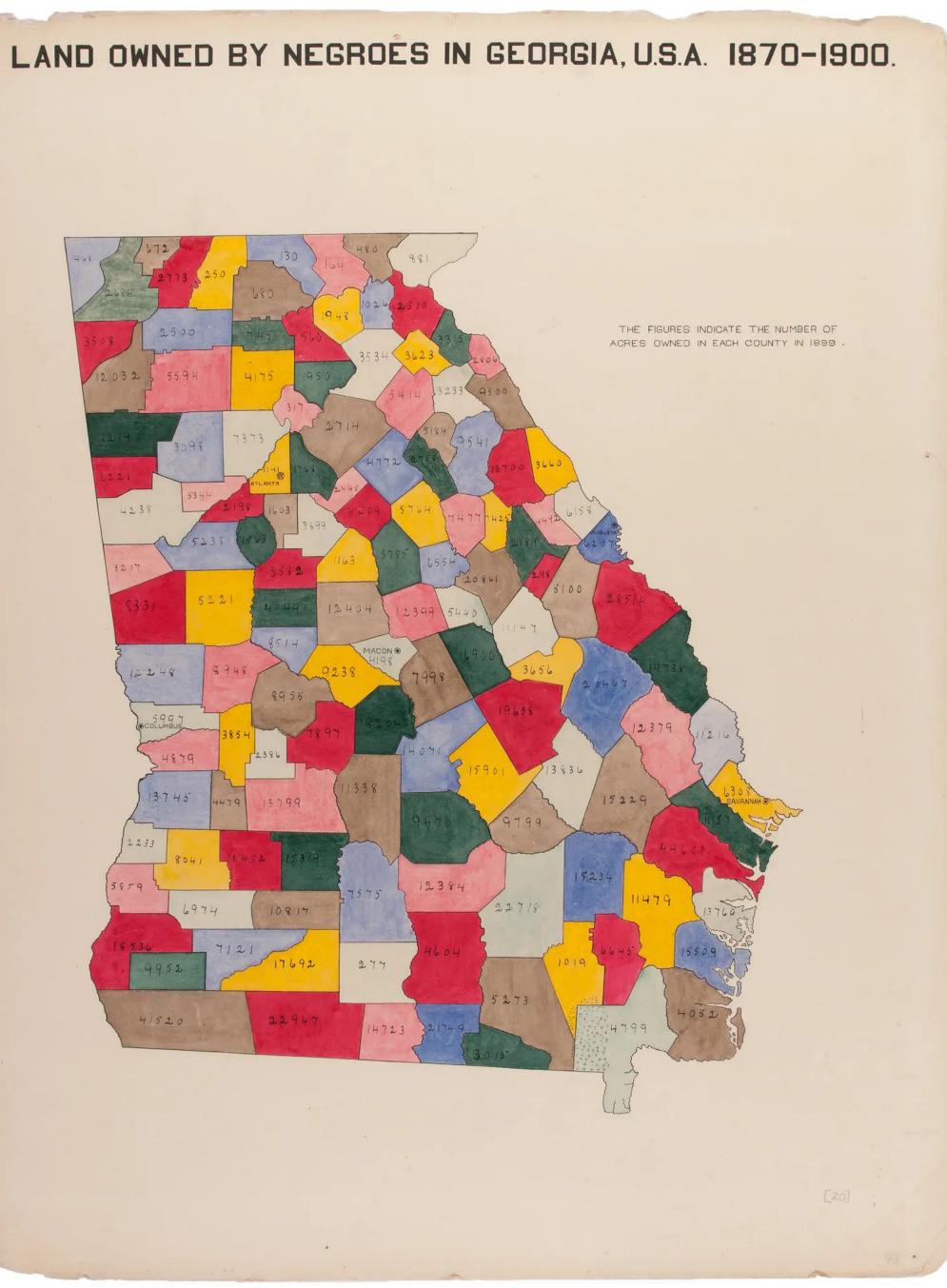
- W. E. B. Dubois (1868-1963) was an African-American historian, sociologist, and activist.
 - He helped create the National Association for the Advancement of Colored People.
 - He is also the first African-American to earn a PhD from Harvard.
- For the 1900 "Exposition Universelle" (world's fair) in Paris, he created a display celebrating the advancements of African-Americans.
 - Part of the display was a series of data visualizations (more in the reading).





https://lithub.com/w-e-b-du-bois-in-paris-the-exhibition-that-shattered-myths-about-black-america/

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Derivation of the Gaussian distribution

Gauss and least squares

- **Recall from Lecture 4:** one of the key differences between the approaches to least squares by Gauss and Legendre was that Gauss linked the theory of least squares to probability theory.
- Specifically, he posed the least squares **model** where

- Where did this distribution come from?

$$y_i = a + bx_i + \epsilon_i$$

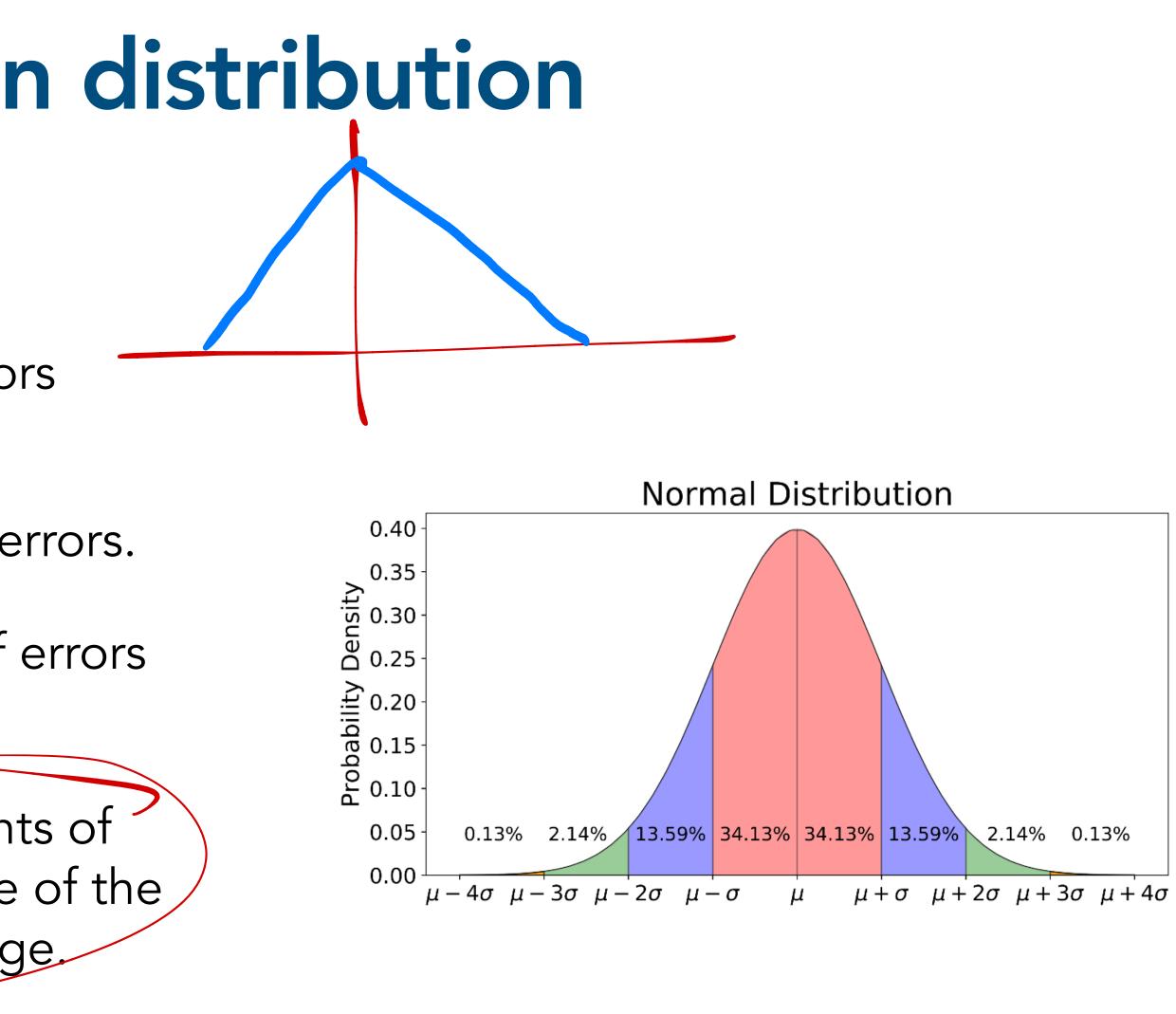
where ϵ_i is a **random variable** that follows the following **error distribution**: $\phi(x \mid \mu, \sigma) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$

Criteria for the Gaussian distribution

- Gauss described that the distribution of errors should satisfy three criteria:
 - Small errors are more likely than large errors.
 - 2. For any real number ϵ the likelihood of errors of magnitude ϵ and $-\epsilon$ are equal.

In the presence of several measurements of the same quantity, the most likely value of the quantity being measured is their average.

We will now derive the Gaussian distribution using just these three criteria. (Buckle up!)



ø(x) "probability density function" (i) $\phi(x)$ maximized when x = 0(2) $\phi(x) = \phi(-x)$ Aside: new function $f(x) = \frac{\phi'(x)}{\phi(x)}$ Poperty of f: f(-x) = -f(x) $f(-x) = \frac{\phi'(-x)}{\varphi(-x)} = \frac{-\phi'(x)}{\varphi(x)} = -f(x)$

"phi of x"

(3) Suppose p is some fixed, unknown gualits that we are trying to measure. => let M1, M2, ..., My be estimates / masurements of p. =) $E_{mor}: M_1 - p, M_2 - p, \dots, M_n - p$ (could be $p - M_i$ too) => Likelihood: $L(p) = \phi(M_1 - p) \cdot \phi(M_2 - p) \cdot \phi(M_3 - p) \cdots \phi(M_n - p)$ Told: the "most likely" p is average $\Rightarrow p = \frac{M_1 + M_2 + \dots + M_n}{n}$ -) L(p) maximized when p= M



 $\cdots \beta(M_n - p)$ $L(p) = \varphi(M, -p) \cdot \varphi(M_2 - p) \cdot \varphi(M_3 - p)$ $\frac{d}{dp}L(q)=0 \quad \text{when} \quad p=\tilde{M}$ $L_{y} = \frac{d}{dr} L(p) = -\phi'(n_{y}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}-p)\phi(n_{z}$ $a'(\pi) = f'(\pi)g(\pi)$ $- \varphi(M, -p) \varphi'(M_2 - p) \varphi(M_3 - p) - ...$ +f(x)g'(x) $- \beta(M, -p) \beta(M_2 - p) \beta'(M_3 - p)$ $= a(x) \cdot \frac{f'(x)}{f(x)} + a(x) \cdot \frac{g'(x)}{g(x)} + \frac{g'(x)}{g(x)}$ $= -L(p) \cdot \frac{\phi'(n,-p)}{\phi(n,-p)} - L(p) \cdot \frac{\phi'(n_2-p)}{\phi(n_2-p)}$



 $\frac{d}{d\rho}L(\rho) = -L(\rho)\cdot\frac{\phi'(n,-\rho)}{\phi(n,-\rho)} - L(\rho)\cdot\frac{\phi'(n_2-\rho)}{\phi(n_2-\rho)} \dots$ $= -L(p) \left[\frac{\hat{S}}{\sum_{i=1}^{\infty} \frac{\varphi'(M_i - p)}{\varphi(M_i - p)} \right]$ $= -\iota(p)\cdot \left| \hat{\xi}f(M_i-p) \right| = 0$ p=M satisfies reed an f that satisfier this $\sum_{i=1}^{n} \mathcal{A}(n_i - \overline{n}) = 0$ for my measurements

 $\forall \Lambda_1, \Lambda_2, \ldots, \Lambda_n$ $\sum_{i=1}^{n} f(n_i - n_i) = 0$ f(kx) = kf(x)

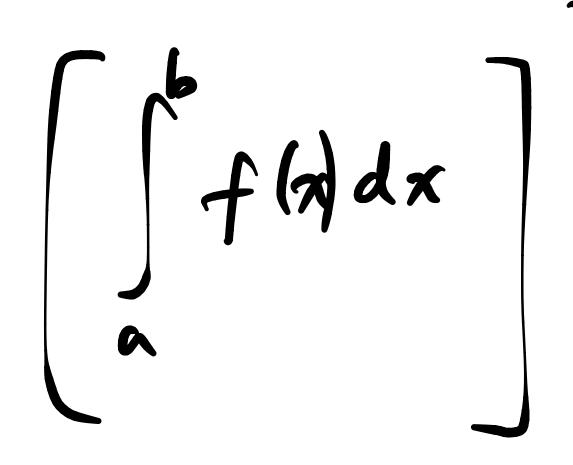
=) key i dea: f must be linear! (linhed reading)

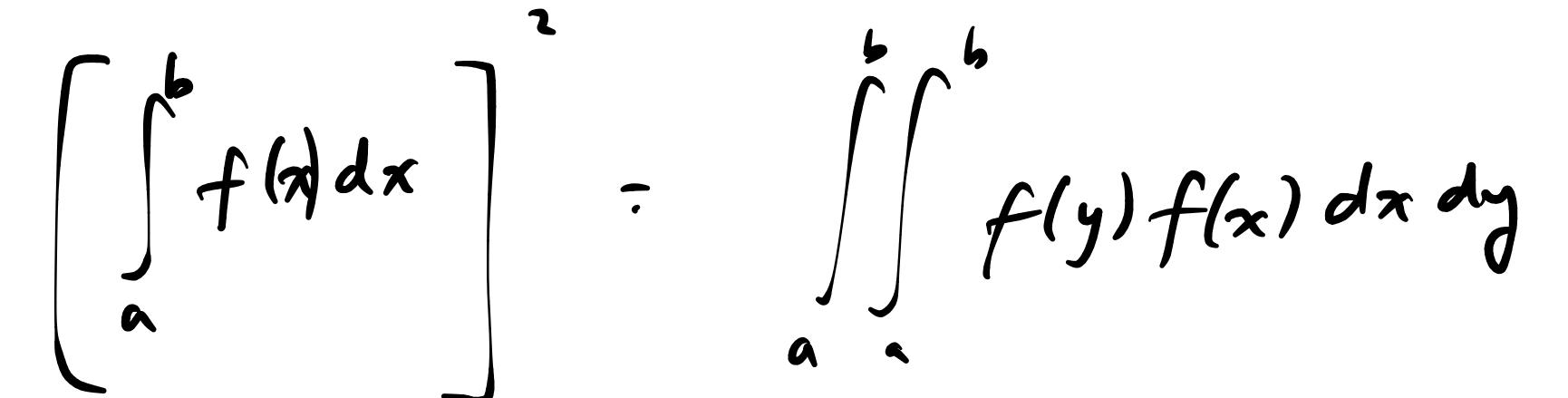
 $\Rightarrow f(x) = CX \qquad C$



f(x) = C x $\int \frac{1}{n} dn = \ln \left[n \right]$ - P $\phi(x) = e^{\frac{1}{2}x}$ $\phi'(x) = Cx$ **6**(x) stagrate both sider $\sigma'(x) = \int dx = \int dx$ pu st $\frac{S}{2}x^{2} + D$ ln ø(x) (\mathbf{l}) $\int e^{-x} dx = \sqrt{TT}$ corditations 2010







Origins of computation

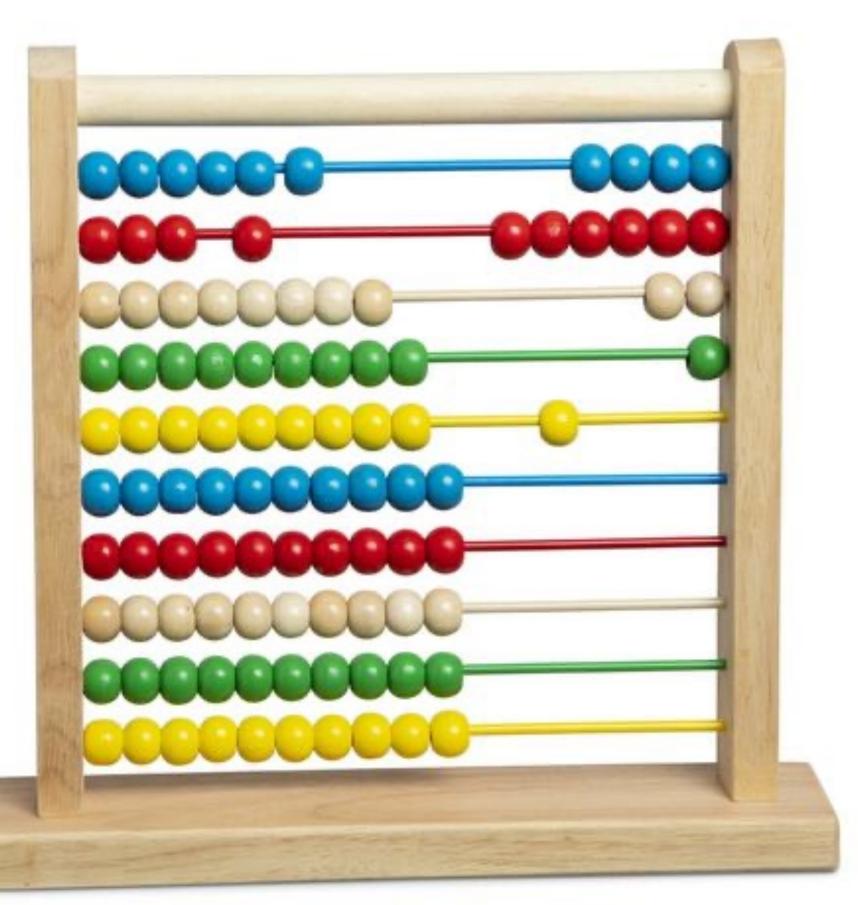






- Abacuses are still in use today!
- It can be used to evaluate addition, subtraction, multiplication, and division without needing symbolic arithmetic.
 - The number system we use today is known as the Hindu-Arabic number system, and dates to ~700 AD.
- See <u>this site</u> for a digital abacus.
- 1. <u>https://criticallyconsciouscomputing.org/history</u>

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Computers

- "computer".
- algebraic operations.
- "Computer"s referred to people as recent as the mid-1900s.
 - men.
- 1. <u>https://criticallyconsciouscomputing.org/history</u>

 In 1613¹, English poet Richard Braithwait published a book titled The Yong Mans Gleanings, which contains the earliest known reference to the word

In it, a "computer" was defined as a person who performed arithmetic and

Then, most computers were women, because they could be paid less than

Leibniz and binary

- course is also credited for developing the **binary number system**.
 - There are two digits in binary 1 and 0. "Bit" stands for "binary digit."
- level in binary.
 - or "off" (0).
- CHINESE FIGURES OF FU XI".

• Gottfried Wilhelm Leibniz (1454-1716) – one of the "creators of calculus" that we studied earlier in the

• At the time, he had no practical use for it, but in modern computing all code is represented at a low-

• Transistors – the building block of computer processors – are like switches that can either be "on" (1)

• One of this week's readings contains a translation of his original work that discussed binary numbers, titled "EXPLANATION OF BINARY ARITHMETIC, WHICH USES ONLY THE CHARACTERS 0 AND 1, WITH SOME REMARKS ON ITS USEFULNESS, AND ON THE LIGHT IT THROWS ON THE ANCIENT





Number systems

- Our standard number system, base 10, has 10 digits: 0, 1, 2, 3, 4, 5, 6, 7, 8, and 9. 3429 • We can decompose the number 3249_{10} as follows:

$$3429_{10} = 3 \cdot 10^3 + 4 \cdot 10$$

- Similarly, base 2 (i.e. binary) only has 2 digits: 0 and 1.
 - To convert from binary to base 10, we can follow a similar procedure.

$$1001_{10} = 1 \cdot 2^3 + 0 \cdot 2^2 + 2^3$$

decimal dec = 10

 $)^2 + 2 \cdot 10^1 + 9 \cdot 10^0$

 $rac{0}{0} \cdot 2^{1} + 1 \cdot 2^{0} = 8 + 1 = 9$

Converting between binary and base 10

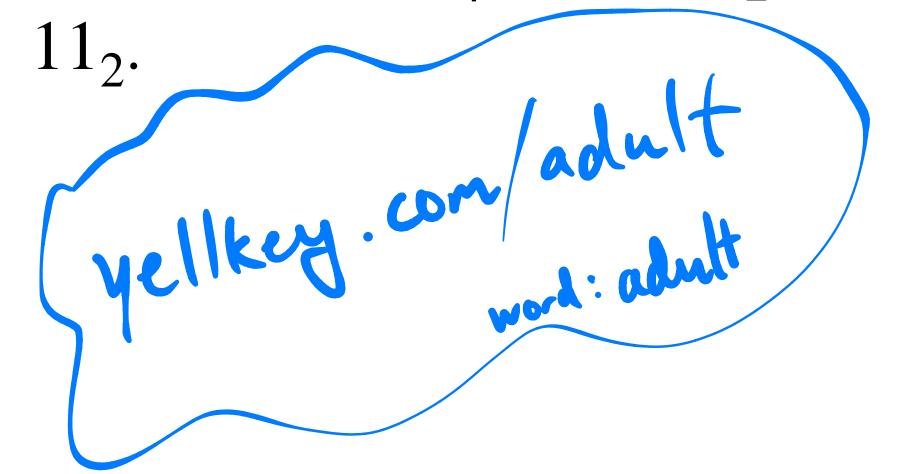
Example: Convert 324_{10} to binary. 324 = 256 + 64 + 4 $31 = 2^{8} + 2^{4} + 2^{2}$ - 256



Example: Convert 110001011_{2} to base 10. $2^{2}+2^{2}+2^{2}+2^{2}+2^{2}+2^{2}=$ = 395

Binary arithmetic

- Arithmetic in binary largely works the same way that arithmetic in base 10 works.
- **Example:** multiply 10110_2 and



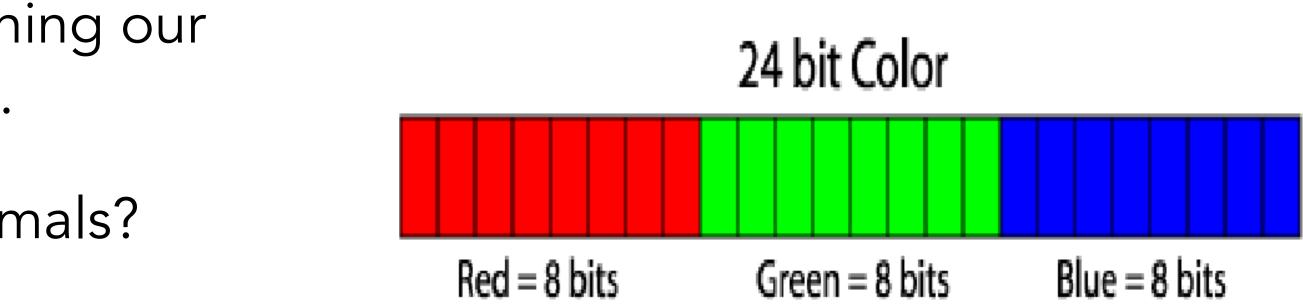


 $- 2^{\circ} + 2^{\prime} + 2^{3} = 11$ (D $10101, -2^{\circ}+2^{2}+2^{4}=121$ (0|10(0) | 0X

000010

Symbols

- We've looked at how base 10 numbers can be stored in binary.
- But base 10 numbers are not the only thing our computers need to store and work with.
 - What about negative numbers? Decimals?
 - Strings?
 - Colors?
 - All of these can be stored in binary as well.



Boolean algebra

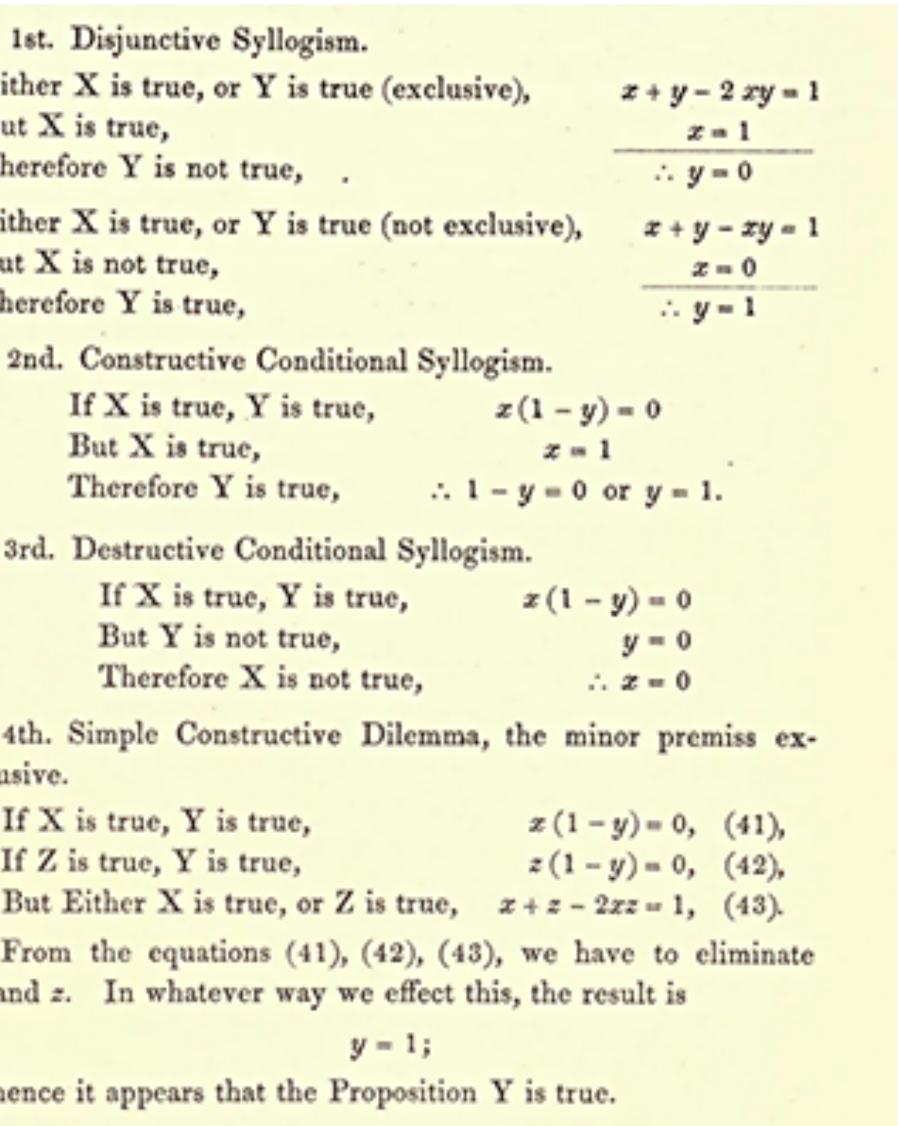
- George Boole (1815-1864) was an English mathematician.
- In 1854, he published An investigation into the Laws of Thought, on Which are founded the Mathematical Theories of Logic and Probabilities, in which he laid the foundations of **Boolean algebra**.
 - In Boolean algebra, there are two values 1 (True) and 0 (False), and three operators – AND, OR, and NOT.
- Fun fact: at the age of 19, Boole created his own elementary school!

| 1st. Disjunctive Syllogism. |
|----------------------------------------------------------------|
| Either X is true, or Y is true (exclusive), $x + y - 2xy =$ |
| But X is true, $x = 1$ |
| Therefore Y is not true, $\therefore y = 0$ |
| Either X is true, or Y is true (not exclusive), $x + y - xy =$ |
| But X is not true, $x = 0$ |
| Therefore Y is true, $\therefore y = 1$ |
| 2nd. Constructive Conditional Syllogism. |
| If X is true, Y is true, $x(1-y) = 0$ |
| But X is true, $x = 1$ |
| Therefore Y is true, $\therefore 1 - y = 0$ or $y = 1$. |
| 3rd. Destructive Conditional Syllogism. |
| If X is true, Y is true, $x(1-y) = 0$ |
| But Y is not true, $y = 0$ |
| Therefore X is not true, $\therefore x = 0$ |
| 4th. Simple Constructive Dilemma, the minor premiss exclusive. |
| If X is true, Y is true, $x(1-y) = 0$, (41), |
| If Z is true, Y is true, $z(1-y) = 0$, (42), |
| But Either X is true, or Z is true, $x + z - 2xz = 1$, (43). |
| From the equations (41) (42) (43) we have to eliminat |

y = 1;

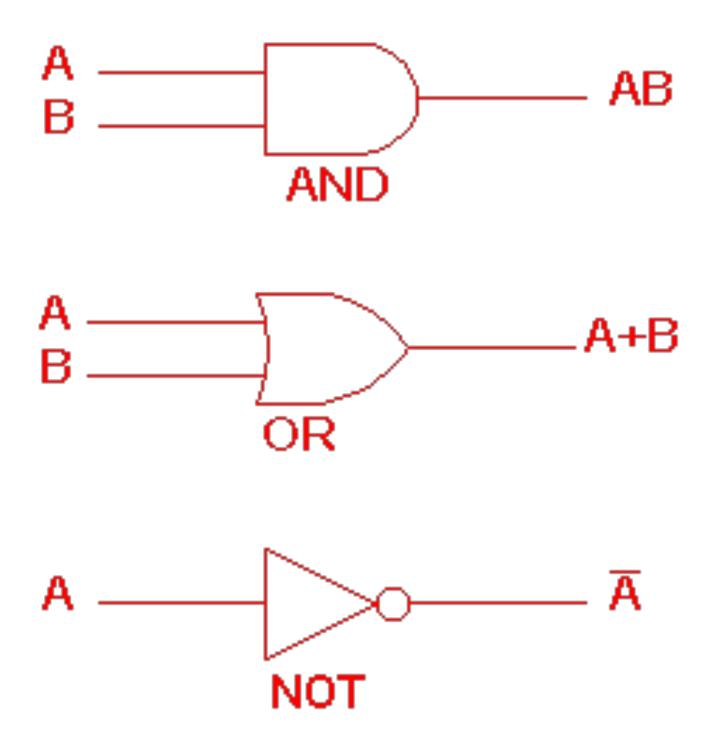
whence it appears that the Proposition Y is true.

x and z. In whatever way we effect this, the result is



Boolean algebra

- In Boolean algebra, there are two values – 1 (True) and 0 (False), and three operators – AND, OR, and NOT.
- All other operations can be constructed using a combination of these three operators.
- Circuits use Boolean algebra to control the flow of current.



That's all!